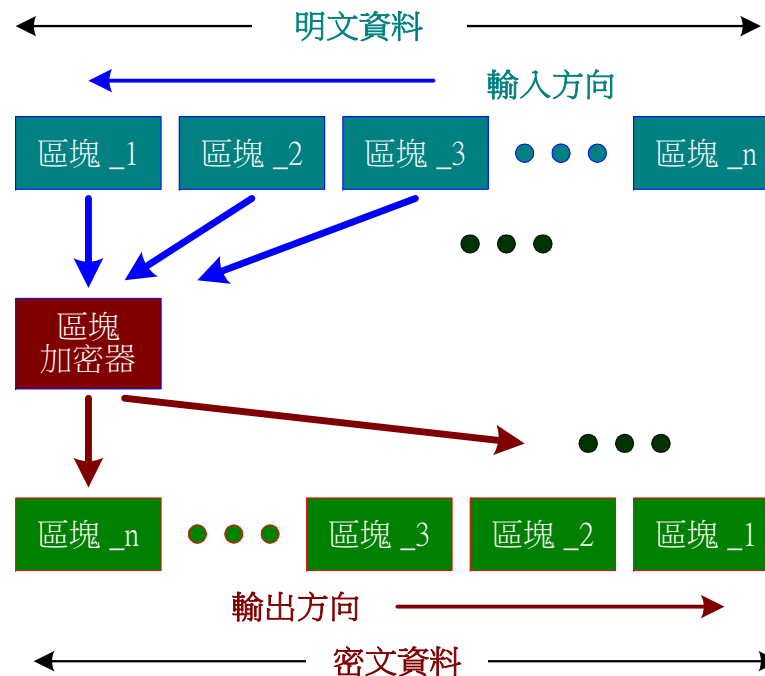


DES 操作模式



- ✦ 電子密碼書 (Electronic Code Book, ECB) 模式
- ✦ 密文區段串接 (Cipher Block Chaining, CBC) 模式
- ✦ k-位元密文反饋 (k-bits Cipher Feedback, CFB) 模式
- ✦ k-位元輸出反饋 (k-bits Output Feedback, OFB) 模式



DES 電子密碼書模式



✦ 電子密碼書模式 (ECB)

◆ 加密處理：

• 明文： $P = P_1 \parallel P_2 \parallel P_3, \dots, \parallel P_N$

• 密文： $C = E_K(P) = E_K(P_1) \parallel E_K(P_2) \parallel \dots \parallel E_K(P_N)$

$E_K(P_N)$

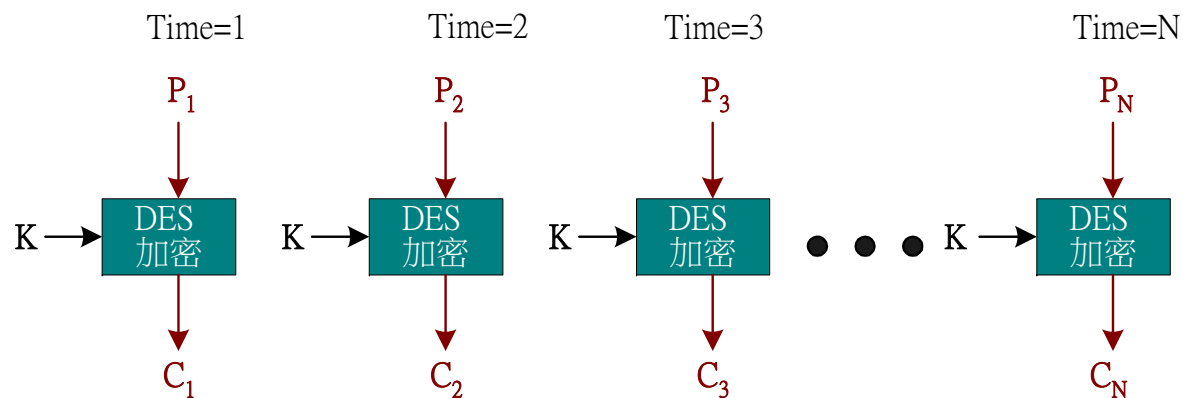
◆ 解密處理：

• 密文： $C = C_1 \parallel C_2 \parallel C_3, \dots, \parallel C_N$

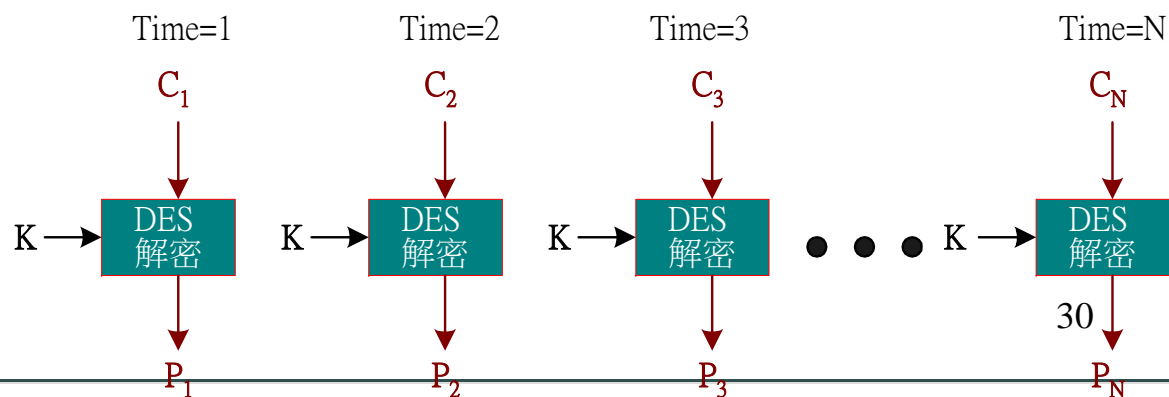
• 明文： $P = E_K(C) = E_K(C_1) \parallel E_K(C_2) \parallel \dots \parallel E_K(C_N)$

◆ 容易遭受明文攻擊

(a) ECB 加密操作 $C = E_K(P)$



(b) ECB 解密操作 $P = D_K(C)$



DES 密文區塊串接模式 (1)



✦ 密文區塊串接模式(Cipher Block Chaining , CBC)

◆ 增加『初始向量』(Initialization Vector, IV)

◆ 加密程序：

- $C_1 = E_K(IV \oplus P_1)$
- $C_i = E_K(C_{i-1} \oplus P_i) ; i = 2, 3, 4, \dots, N$
- $C = C_1 \parallel C_2 \parallel C_3, \dots, C_N$

◆ 解密程序：

- $P_1 = D_K(C_1) \oplus IV$
- $P_i = D_K(C_i) \oplus C_{i-1} ; i = 2, 3, 4, \dots, N$
- $P = P_1 \parallel P_2 \parallel P_3, \dots, P_N$

◆ 驗證：

- $D_K(C_i) = D_K(E_K(C_{i-1} \oplus P_i)) ; i = 2, 3, 4, \dots, N$
 $= (C_{i-1} \oplus P_i)$

• 則： $D_K(C_i) \oplus C_{i-1} = (C_{i-1} \oplus P_i) \oplus C_{i-1} = P_i$ 得證之。

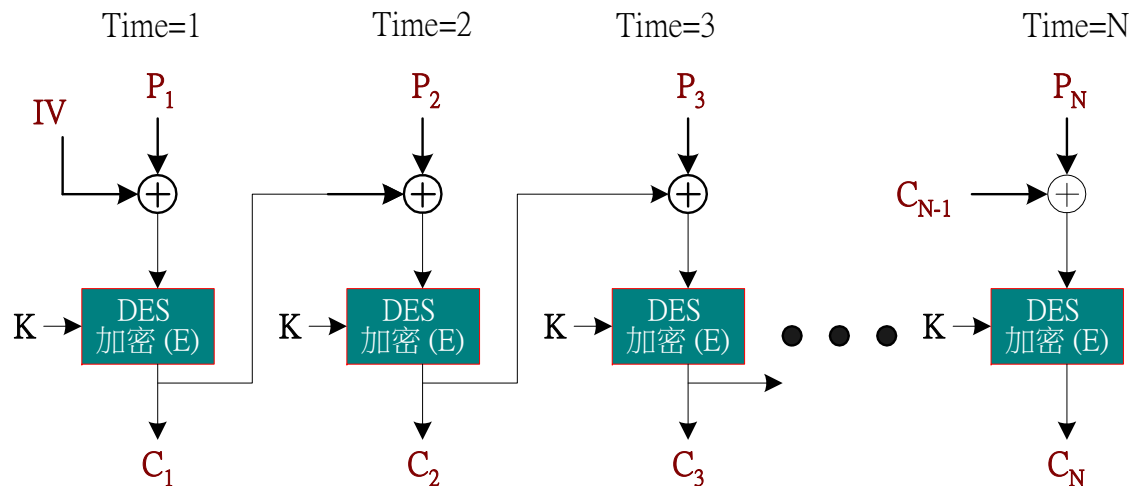


DES 密文區塊串接模式 (2)

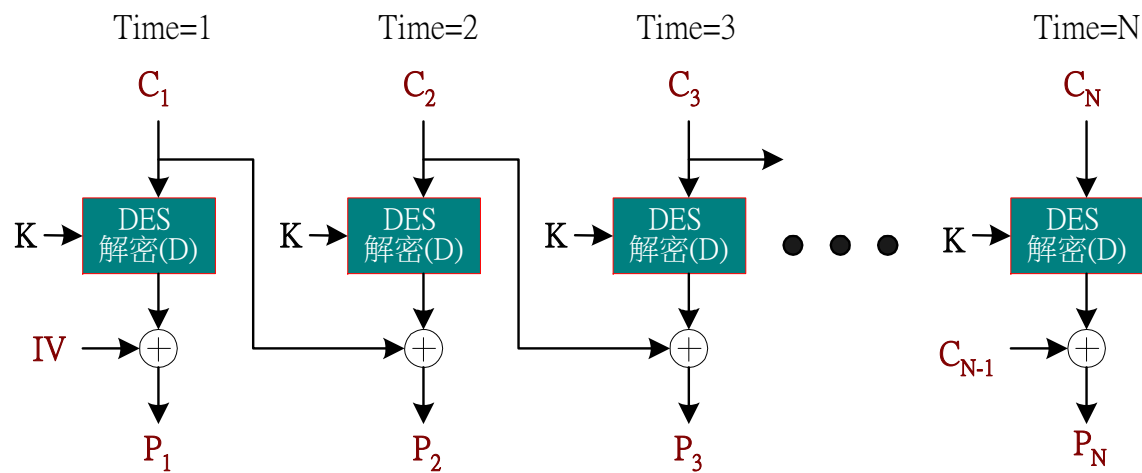


☀ CBC 運作程序

(a) CBC 加密操作 $C = E_K(P)$



(b) CBC 解密操作 $P = D_K(C)$



DES J-位元密文反饋模式(1)



* J-位元密文反饋模式(J-bits Cipher Feedback , CFB)

◆ J 位元串流加密 (Stream Cipher)

◆ 加密運算程序：

- $SR_1 = IV$
- $C_1 = F_j(E_K(SR_1)) \oplus P_1$
- $SR_m = S_j(SR_{m-1}) \parallel C_{m-1} ; m = 2, 3, 4, \dots, N$
- $C_m = F_j(E_K(SR_m)) \oplus P_m ; m = 2, 3, 4, \dots, N$
- $C = C_1 \parallel C_2 \parallel C_3, \dots, C_N$

◆ 解密運算程序：

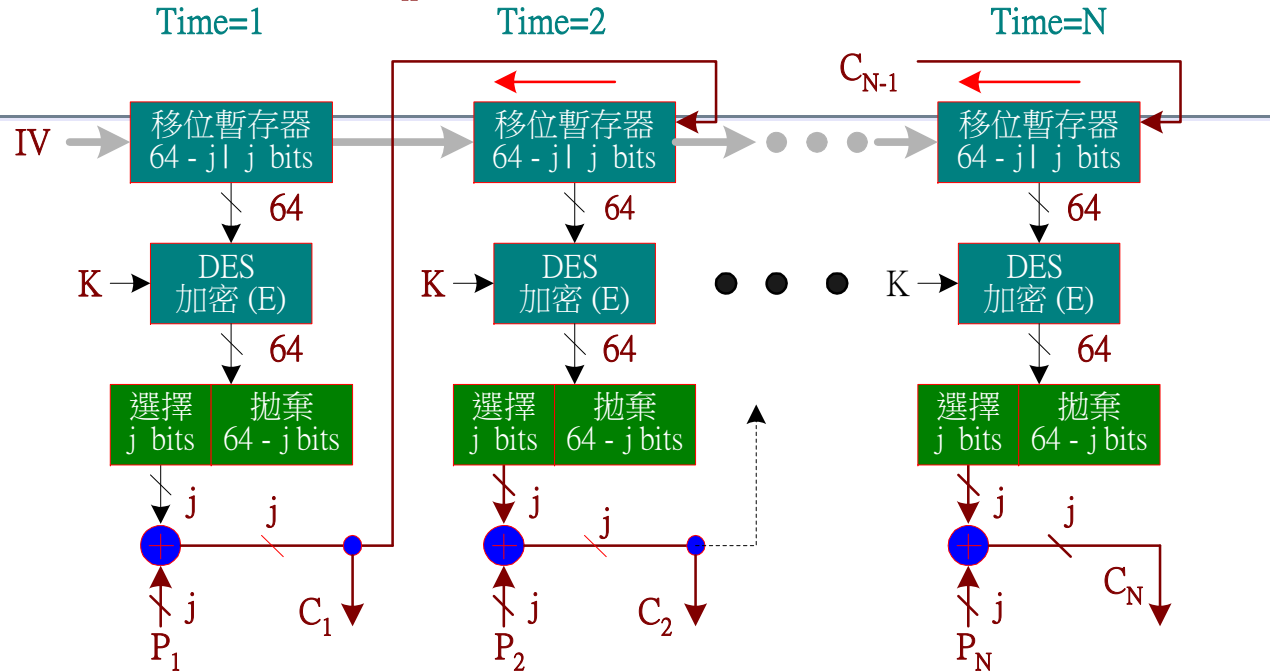
- $SR_1 = IV$
- $P_1 = F_j(D_K(SR_1)) \oplus C_1$
- $SR_m = S_j(SR_{m-1}) \parallel C_{m-1} ; m = 2, 3, 4, \dots, N$
- $P_m = F_j(D_K(SR_m)) \oplus C_m ; m = 2, 3, 4, \dots, N$
- $P = P_1 \parallel P_2 \parallel P_3, \dots, P_N$



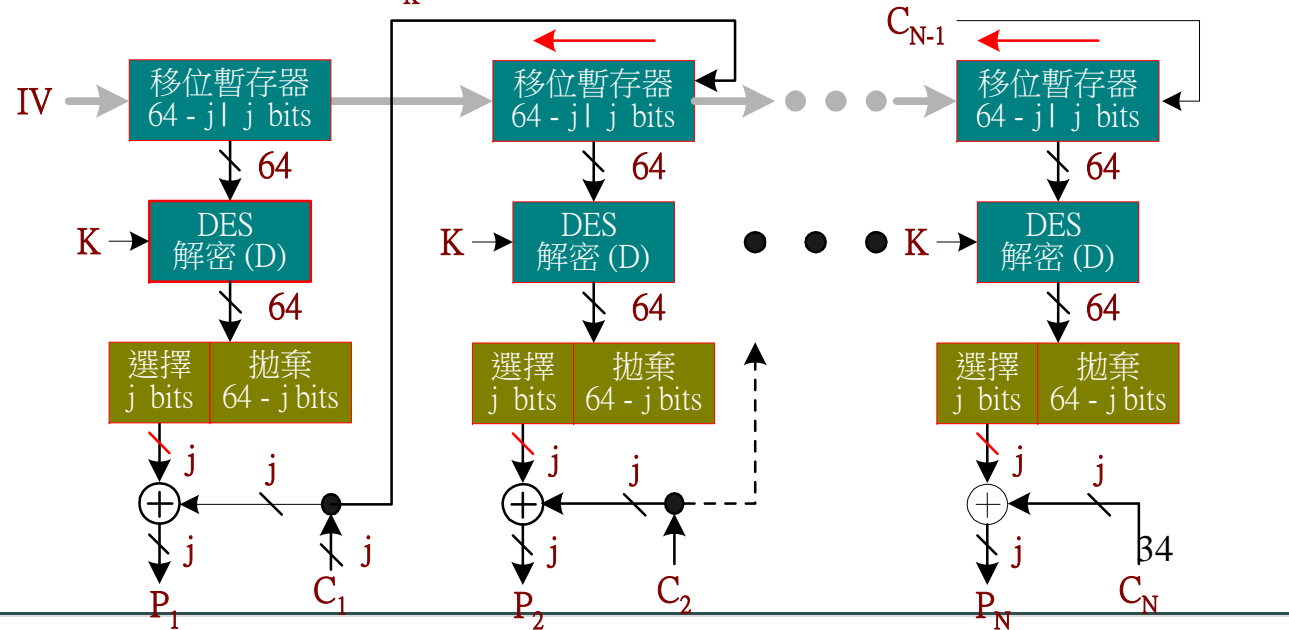
DES J-位元 密文反饋模 式 (2)

CFB 運作程序

(a) CFB 加密操作 $C = E_K(P)$



(b) CFB 解密操作 $P = D_K(C)$



DES J-位元輸出反饋模式(1)



✦ J-位元輸出反饋模式(J-bits Output Feedback, OFB)

◆ J-位元串流加密

◆ 加密運算乘程序：

- $SR_1 = IV$
- $O_1 = F_j(E_K(SR_1))$
- $C_1 = P_1 \oplus O_1$
- $SR_m = S_j(SR_{m-1}) \parallel O_{m-1} ; m = 2, 3, 4, \dots, N$
- $O_m = F_j(E_K(SR_m)) ; m = 2, 3, 4, \dots, N$
- $C_m = O_m \oplus P_m ; m = 2, 3, 4, \dots, N$
- $C = C_1 \parallel C_2 \parallel C_3, \dots, C_N$

◆ 解密運算乘程序：

- $SR_1 = IV$
- $O_1 = F_j(D_K(SR_1))$
- $P_1 = O_1 \oplus C_1$
- $SR_m = S_j(SR_{m-1}) \parallel O_{m-1} ; m = 2, 3, 4, \dots, N$
- $O_m = F_j(D_K(SR_m)) ; m = 2, 3, 4, \dots, N$
- $P_m = O_m \oplus C_m ; m = 2, 3, 4, \dots, N$
- $P = P_1 \parallel P_2 \parallel P_3, \dots, P_N$

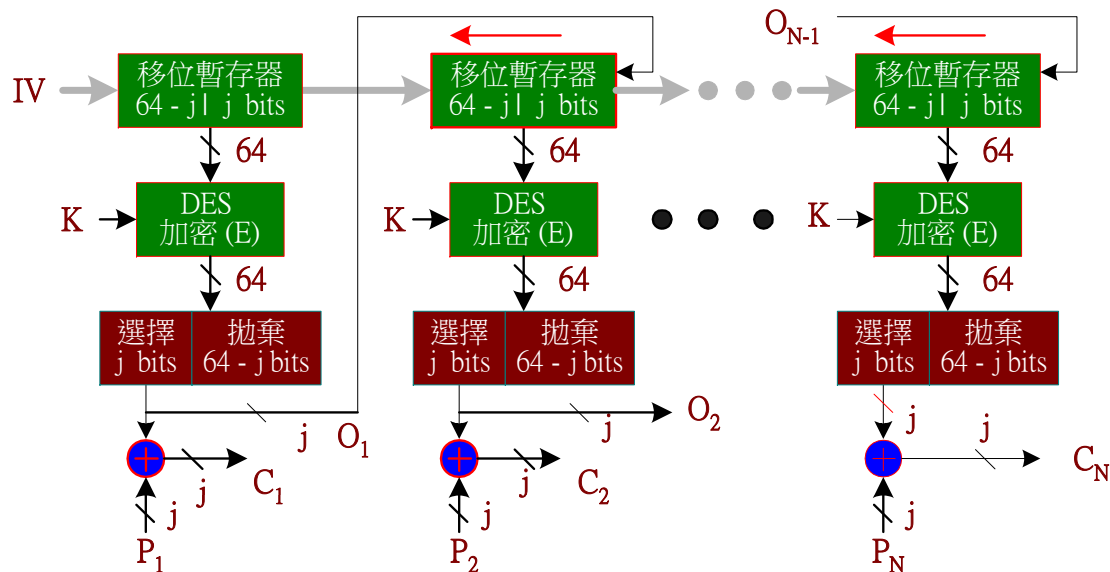




DES J-位元 輸出反饋模 式 (2)

☀ OFB 運作程序

(a) OFB 加密操作 $C = E_K(P)$



(b) OFB 解密操作 $P = D_K(C)$

